

Govt. Engg. College Jhalawar

Sub:- TOC

MM- 10

Hr.-01

- Q1. What is automation? Differentiate the DFA and N DFA. (3)
- Q2. Prove that for any transition function δ and for any two input strings x and y , $\delta(q,xy) = \delta(\delta(q,x),y)$. (2)
- Q3. Construct a grammar G accepting the set L of all strings over $\{ a,b \}$ having more a 's than b 's. (2)
- Q4. Construct a nondeterministic finite automaton accepting $\{ab,ba\}$, and use it to find a deterministic automaton accepting the same set. (3)

Subject :- TOC

Q.1
Ans

What is automation? Differentiate the DFA and N DFA.
An Automation is defined as a system where energy, materials and information are transformed, transmitted and used for performing some operation/ functions without direct participation of man. Examples are automation is automatic machine tools, automatic parking machines etc.

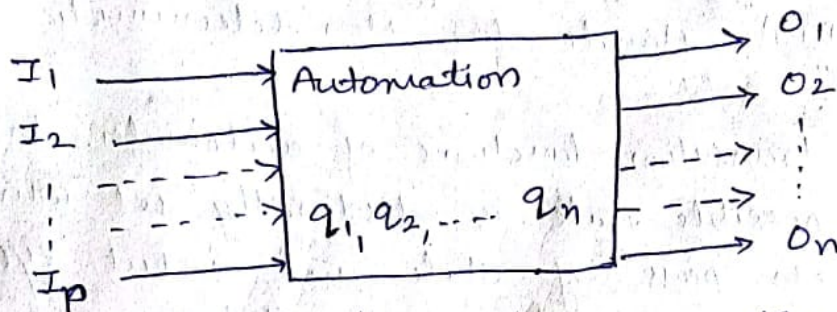


Fig:- Model of discrete automation.

The characteristics of automation are now described.

- (i) I/P (Input) :- At each of the discrete instants of time t_1, t_2, \dots, t_n the input values I_1, I_2, \dots, I_p , each of which can take a finite number of fixed values from the input alphabet Σ , are applied to the input side of the model.
- (ii) output (O/P) :- O_1, O_2, \dots, O_n are the output of the model, each of which can take a finite number of fixed values from an output O .
- (iii) State :- At any instant of time the automation can be in one of the state q_1, q_2, \dots, q_n .
- (iv) State Relation :- The next state of an automation at any time is determined by the present state and the present Input.

(v) output relation:- The output is related to the either state only or to both the input and the state. It should be noted that at any instant of time the automation is in some state. On reading an input symbol, the automation moves to a next state which is given by the state relation.

Difference b/w DFA and NFA.

- ① "DFA" stands for "Deterministic finite automata" while "NFA" stands for "Non deterministic finite automata".
- ② Both the transition functions of automata. In DFA the next possible state is distinctly set while in NFA each pair of state and input symbol can have many possible next states.
- ③ DFA cannot use empty string state while NFA can use empty string transition.
- ④ NFA is easier to construct while it is more difficult to construct DFA.
- ⑤ Backtracking is allowing in DFA while in NFA it may or may not be allowed.
- ⑥ DFA requires more space while NFA requires less space.
- ⑦ While DFA can be understood one machine and a DFA machine can be constructed for every I/P and O/P.
- ⑧ NFA can be understood as ~~one~~ ^{several} machines ~~and a DFA machine~~ can be little machines that computer together, and there is no possibility of constructing an NFA machine for every input and output.

Q2 prove that for any transition function δ and for any two input string x and y , $\delta(q, xy) = \delta(\delta(q, x), y)$.

Ans By the method of Induction, on $|y|$, i.e. length of y .

Basis: when $|y| = 1$, $y = a \in \Sigma$

$$\begin{aligned} \text{L.H.S} &= \delta(q, xa) \\ &= \delta(\delta(q, x), a) \\ &= \text{R.H.S.} \end{aligned}$$

Assume the result for all strings x and string y with $|y| = n$. Let y be a string of length $n+1$ write $y = y_1 a$ where $|y_1| = n$.

$$\begin{aligned} \text{L.H.S} &= \delta(q, xy_1 a) = \delta(q, x_1 a), \quad x_1 = xy_1 \\ &= \delta(\delta(q, x_1), a) \\ &= \delta(\delta(q, xy_1), a) \\ &= \delta(\delta(\delta(q, x), y_1), a) \end{aligned}$$

$$\begin{aligned} \text{R.H.S} &= \delta(\delta(q, x), y_1 a) \\ &= \delta(\delta(\delta(q, x), y_1), a) \end{aligned}$$

L.H.S = R.H.S hence proved.

Q3 Construct a grammar G accepting the set L of all strings over $\{a, b\}$ having more a's ^{than} b's.

Ans for generating strings with a's but with no b's.

$$S \rightarrow a \mid aS \mid Sa$$

for generating strings with a's and b's (more a than b's)

$$S \rightarrow bSS \mid SbS \mid SSb.$$

So, G as follow
 $G = (\{S\}, \{a, b\}, P, S)$ where P consists of

$$S \rightarrow a \mid aS \mid Sa \mid bSS \mid SbS \mid SSb$$

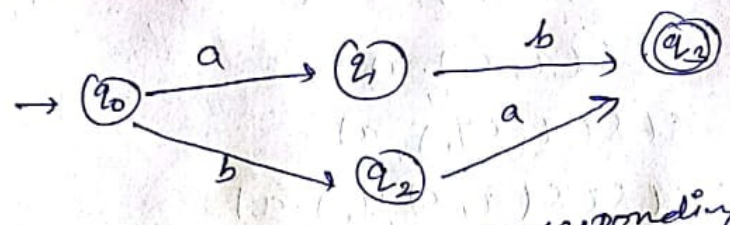
G accept all the strings over $\{a, b\}$ having more a's than b's.

Q4: Construct a N DFA accepting $\{ab, ba\}$ and use it to find a deterministic automation (DFA) accepting the same set.

Ans: The state of N DFA accepting $\{ab, ba\}$ is defined by the table

state/ ϵ	a	b
q_0	q_1	q_2
q_1	-	q_3
q_2	q_3	-
q_3	-	-

The transition diagram of N DFA



The state table of the corresponding DFA is defined by table

state	a	b
$[q_0]$	$[q_1]$	$[q_2]$
$[q_1]$	ϕ	$[q_3]$
$[q_2]$	$[q_3]$	ϕ
$[q_3]$	ϕ	ϕ
ϕ	ϕ	ϕ

The transition diagram of DFA

